

Computational Geometry Algorithms Library

Pierre Alliez INRIA

Andreas Fabri GeometryFactory

http://www.cgal.org/siggraph2009

The updated handout has

- extensive comments in the "Notes" part of Powerpoint.
- hyperlinks to the CGAL User and Reference Manual (~3500 pages).
- hyperlinks to precompiled demos illustrating the algorithms. They will be made available online.

Course Outline

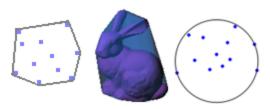
- General Introduction
- CGAL for 2D Vector Graphics
- CGAL for Point Sets
- CGAL for Modeling and Processing of Polyhedral Surfaces
- CGAL for Mesh Generation
- Questions and Answers

Mission Statement

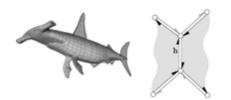
"Make the large body of geometric algorithms developed in the field of computational geometry available for industrial applications"

CGAL Project Proposal, 1996

Algorithms and Datastructures

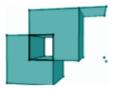


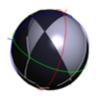
Bounding Volumes



Polyhedral Surface

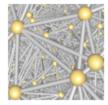


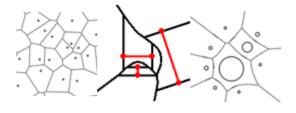




Boolean Operations

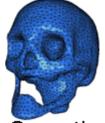


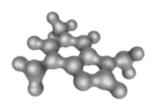




Voronoi Diagrams







Triangulations



Mesh Generation







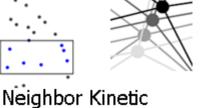


Parameterization Streamlines









Subdivision Simplification



Polytope



Lower Envelope

Arrangement

Intersection Minkowski Detection

Sum

distance

QP Solver

CGAL in Numbers

```
500,000
          lines of C++ code
 10,000
          downloads/year (+ Linux distributions)
  3,500
          manual pages
  3,000
          subscribers to cgal-announce
  1,000
          subscribers to cgal-discuss
    120
          packages
          commercial users
     90
     20
          active developers
     12
          months release cycle
          licenses: Open Source and commercial
```

Some Commercial Users



















cādence™

























Image









Geosystems



































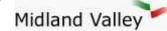










































Why They Use CGAL

"I recommended to the senior management that we start a policy of buying-in as much functionality as possible to reduce the quantity of code that our development team would have to maintain.

This means that we can concentrate on the application layer and concentrate on our own problem domain."

Senior Development Engineer & Structural Geologist

Midland Valley Exploration

Why They Use CGAL

"My research group JYAMITI at the Ohio State University uses CGAL because it provides an efficient and robust code for Delaunay triangulations and other primitive geometric predicates. Delaunay triangulation is the building block for many of the shape related computations that we do. [...]

Without the robust and efficient codes of CGAL, these codes could not have been developed. "

Tamal Dey Professor, Ohio State University

CGAL Open Source Project

Project = « Planned Undertaking »

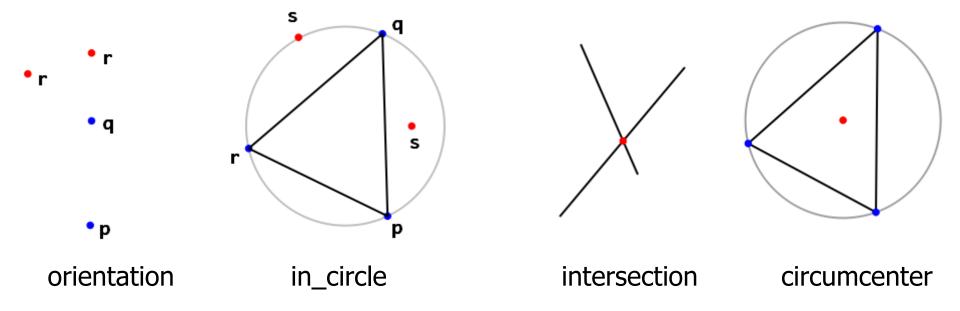
- Institutional members make a long term commitment: Inria, MPI, Tel-Aviv U, Utrecht U, Groningen U, ETHZ, GeometryFactory, FU Berlin, Forth, U Athens
- Editorial Board
 - Steers and animates the project
 - Reviews submissions
- Development Infrastructure
 - Gforge: svn, tracker, nightly testsuite,...
 - 120p developer manual and mailing list
 - Two 1-week developer meetings per year

Contributions

- Submission of specifications of new contributions
- Review and decision by the Editorial Board
- Value for contributor
 - Integration in the CGAL community
 - -Gain visibility in a mature project
 - Publication value for accepted contributions

Exact Geometric Computing

Predicates and Constructions

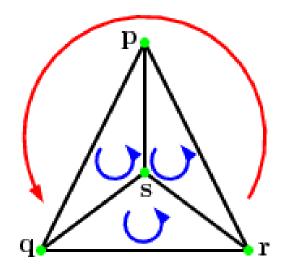


Robustness Issues

- Naive use of floating-point arithmetic causes geometric algorithms to:
 - Produce [slightly] wrong output
 - Crash after invariant violation
 - Infinite loop
- There is a gap between
 - Geometry in theory
 - Geometry with floating-point arithmetic

Geometry in Theory

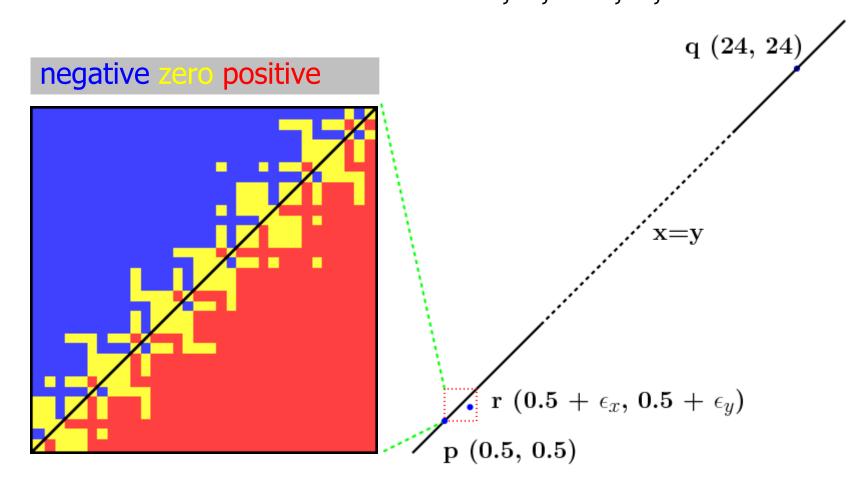
ccw(s,q,r) & ccw(p,s,r) & $ccw(p,q,s) \Rightarrow ccw(p,q,r)$



Correctness proofs of algorithms rely on such theorems

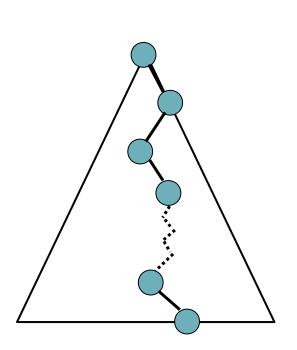
Demo: The Trouble with Double

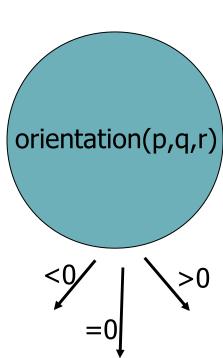
orientation(p,q,r) = sign((p_x-r_x)(q_y-r_y)-(p_y-r_y)(q_x-r_x))



Exact Geometric Computing [Yap]

Make sure that the control flow in the implementation corresponds to the control flow with exact real arithmetic





Filtered Predicates

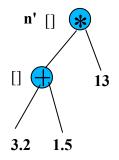
- Generic functor adaptor Filtered_predicate<>
 - Try the predicate instantiated with intervals
 - In case of uncertainty, evaluate the predicate with multiple precision arithmetic

- Refinements:
 - Static error analysis
 - Progressively increase precision

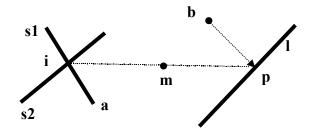
Filtered Constructions

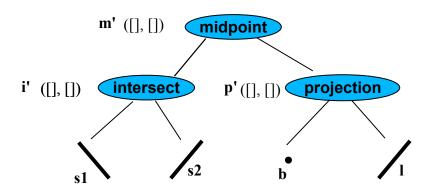
Lazy number = interval and arithmetic expression tree

$$(3.2 + 1.5) * 13$$



Lazy object = approximated object and geometric operation tree





Test that may trigger an exact re-evaluation:

if
$$(n' < m')$$

if (collinear(a',m',b'))

The User Perspective

Convenience Kernels

- Exact_predicates_inexact_constructions_kernel
- Exact predicates exact constructions kernel
- Exact_predicates_exact_constructions_kernel_with_sqrt

Number Types

- double, float
- CGAL::Gmpq (rational), Core (algebraic)
- CGAL::Lazy_exact_nt<ExactNT>

Kernels

- CGAL::Cartesian<NT>
- CGAL::Filtered_kernel<Kernel>
- CGAL::Lazy kernel<NT>

Merits and Limitations

Ultimate robustness inside the black box

- The time penalty is reasonable, e.g. 10% for 3D Delauny triangulation of 1M random points
- Limitations of Exact Geometric Computing
 - Topology preserving rounding is non-trivial
 - Construction depth must be reasonable
 - Cannot handle trigonometric functions

Generic Programming

STL Genericity

```
template <class Key, class Less>
class set {
  Less less;
  insert(Key k)
     if (less(k, treenode.key))
       insertLeft(k);
     else
       insertRight(k);
```

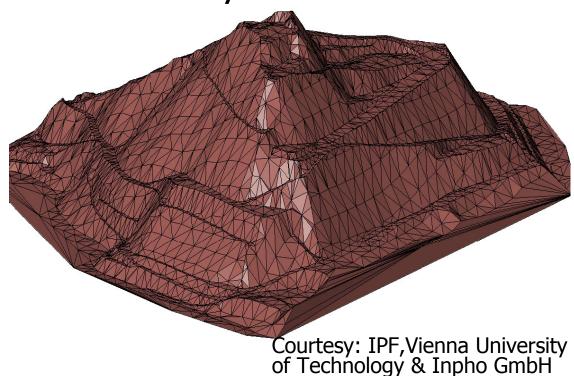
CGAL Genericity

```
template < class Geometry >
class Delaunay triangulation 2 {
    Geometry::Orientation orientation;
    Geometry:: In circle in circle;
    void insert(Geometry::Point t) {
       if(in circle(p,q,r,t)) { . . . }
       if (orientation(p,q,r){...}
```

CGAL Genericity Demo

Without explicit conversion to points in the plane

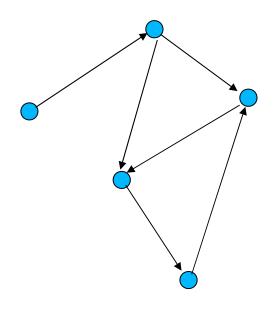
- Triangulate the terrain in an xy-plane
- Triangulate the faces of a Polyhedron



Boost Graph Library (BGL)

- Rich collection of graph algorithms: shortest paths, minimum spanning tree, flow, etc.
- Design that
 - decouples data structure from algorithm
 - links them through a thin glue layer
- BGL and CGAL
 - Provide glue layer for CGAL
 - Extension to embedded graphs inducing the notion of faces

BGL Glue Layer: Traits Class



```
template <typename Graph >
struct boost::graph_traits {
    typedef ... vertex_descriptor;
    typedef ... edge_descriptor;
    typedef ... vertex_iterator;
    typedef ... out_edge_iterator;
};
```

BGL Glue Layer: Free Functions

```
vertex_descriptor v, w;
edge_descriptor e;

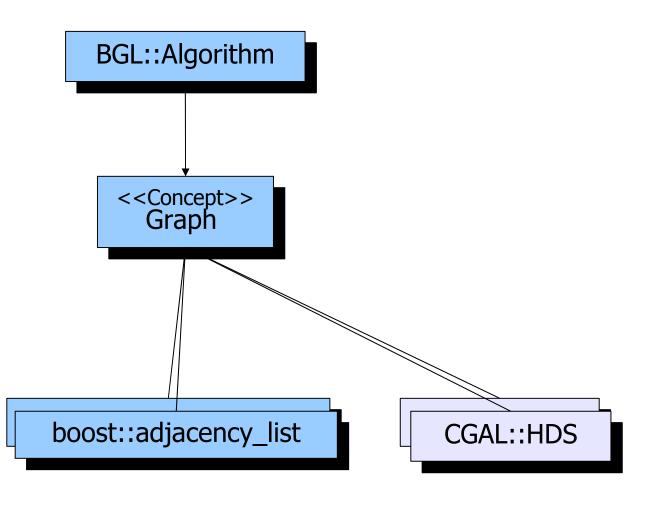
v = source(e,G);
w = target(e,G);
std::pair<out_edge_iterator, out_edge_iterator> ipair;
ipair = out_edges(v,G);
```

BGL Glue Layer for CGAL

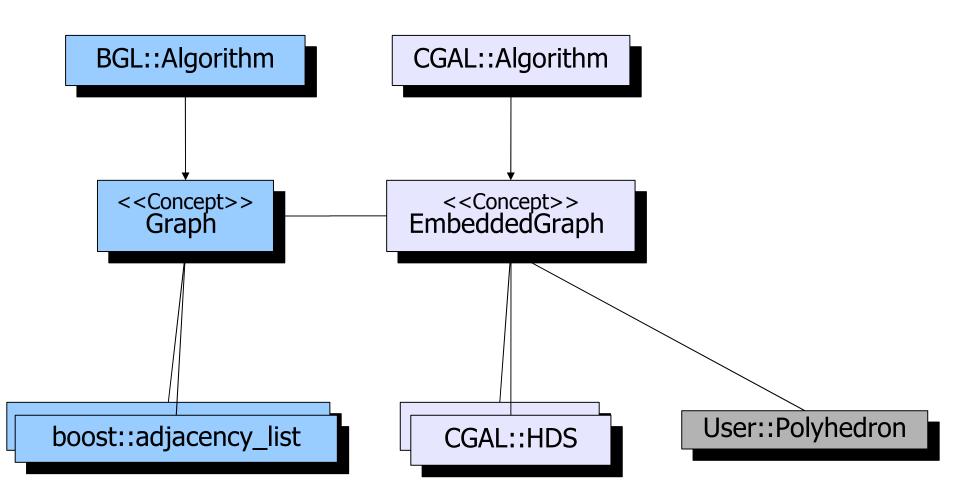
```
CGAL provides partial specializations:
template <typename T>
graph_traits<Polyhedron<T>>;
template <typename T>
Polyhedron<T>::Vertex
source(Polyhedron<T>::Edge);
Users can run:
boost::kruskal mst(P);
```



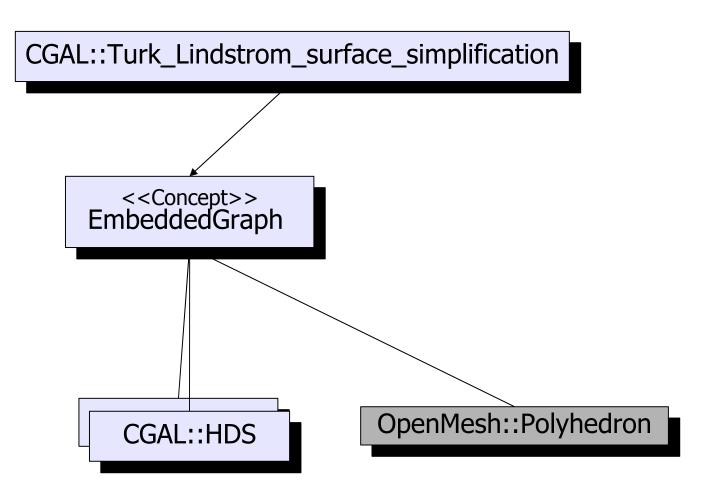
From A BGL Glue Layer for CGAL



To BGL Style CGAL Algorithms



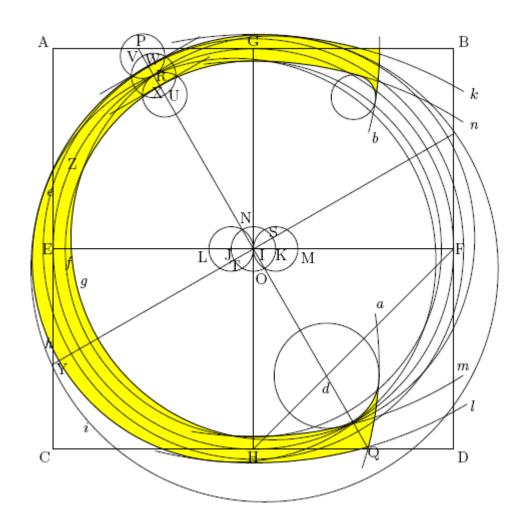
Demo CGAL + OpenMesh



Summary: Overview

Open Source project

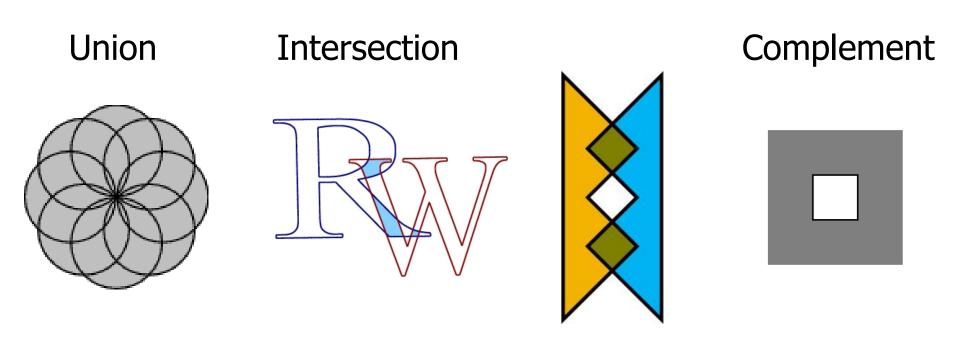
- Clear focus on geometry
- Interfaces with de facto standards/leaders:
 STL, Boost, GMP, Qt, blas
- Robust and fast through exact geometric computing
- Easy to integrate through generic programming



CGAL for 2D Vector Graphics

Andreas Fabri GeometryFactory

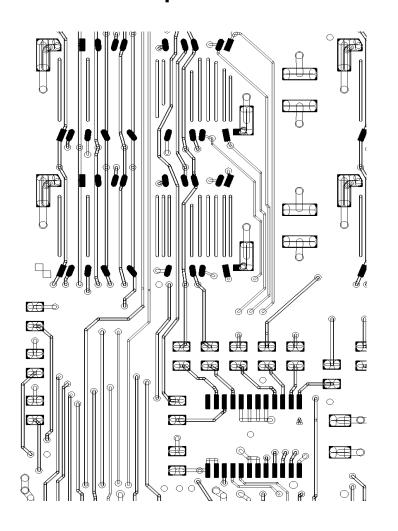
Boolean Operations

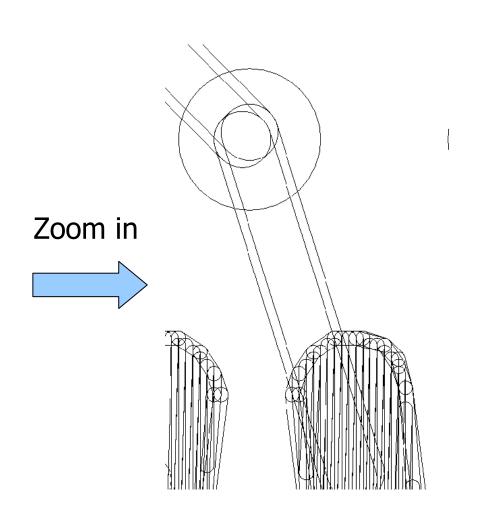


CGAL Boolean Operations can deal explicitly with Circular arcs Bézier Curves Line Segments

Boolean Ops on Circular Arcs Demo

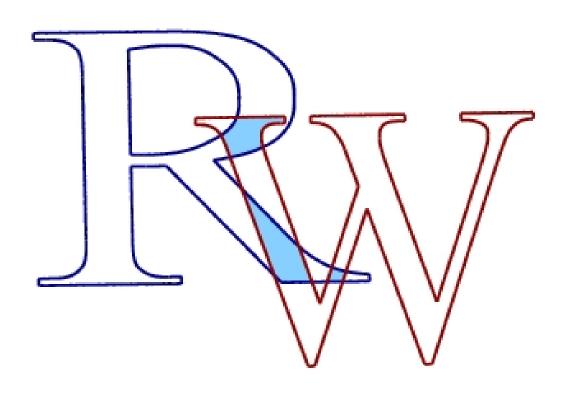
dxf file of a printed circuit board with circular arcs





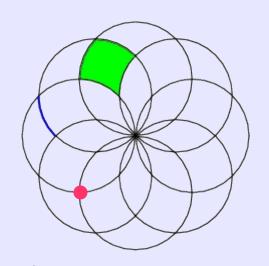
Boolean Ops on Bézier Curves

True Type fonts are Bézier curves

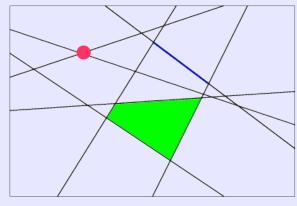


Background: 2D Arrangement

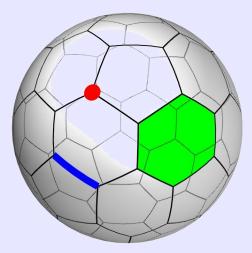
Given a collection of curves on a surface, the **arrangement** is the partition of the surface into **vertices**, **edges** and **faces** induced by the curves



An arrangement of circles in the plane



An arrangement of lines in the plane



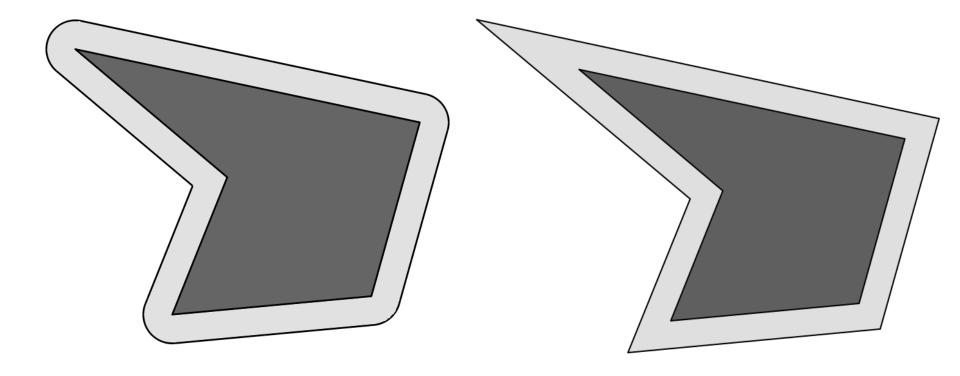
An arrangement of geodesic arcs on the sphere

Arrangement_2<Geometry>

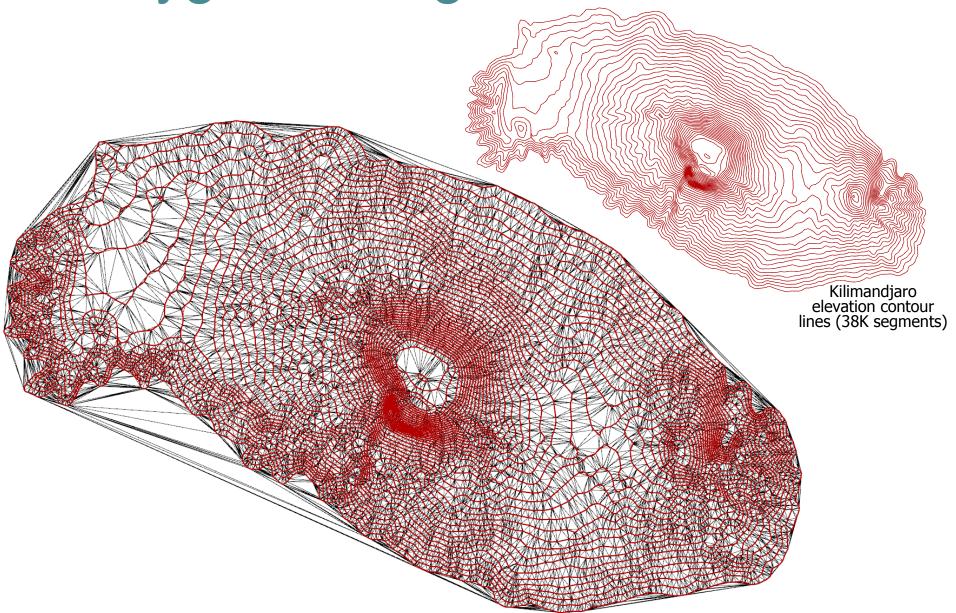
- Constructs, maintains, modifies, traverses, queries, and presents subdivisions of the plane
- Robust and exact
 - All inputs are handled correctly (including degenerate)
 - Exact number types are used to achieve exact results
- Efficient
- Generic
 - Easy to interface, extend, and adapt
 - Notification mechanism for change propagation
- Modular
 - Geometric and topological aspects are separated

Polygon Offsets

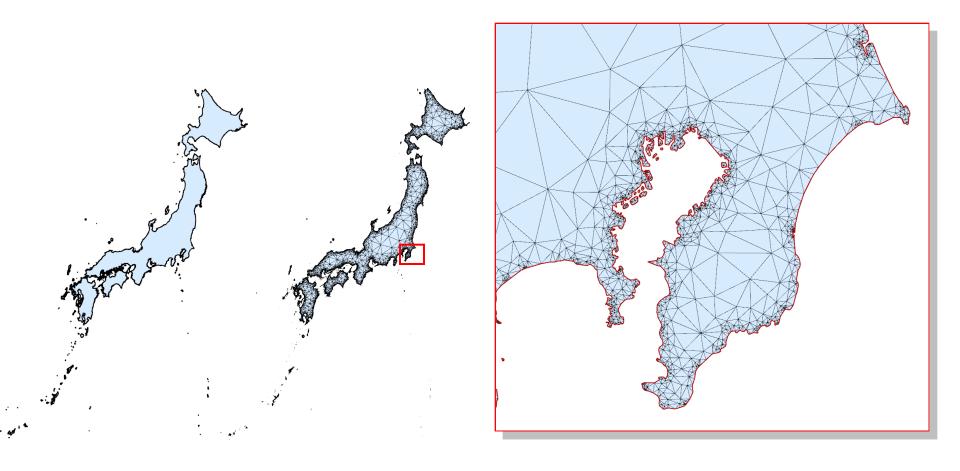
- Based on Minkowski sums, with segments and circular arcs.
- Based on straight skeleton, with segments only.



Polygon Triangulation Demo

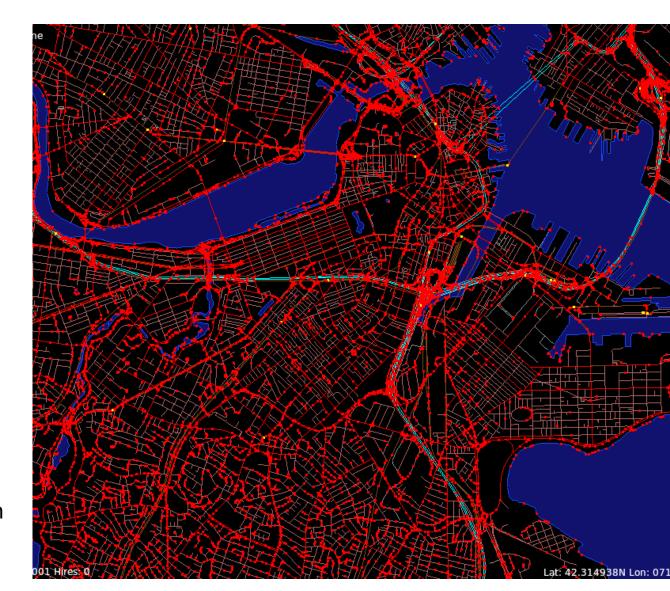


Polygon Mesh Generation Demo



Simultaneous Polyline Simplification

Red points were removed



Courtesy: Laminar Research

Simultaneous Polyline Simplification

Input is the transportation and water layers of OpenStreetMap

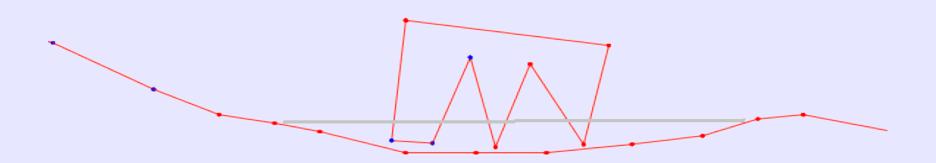
Red points were removed

Courtesy: Laminar Research



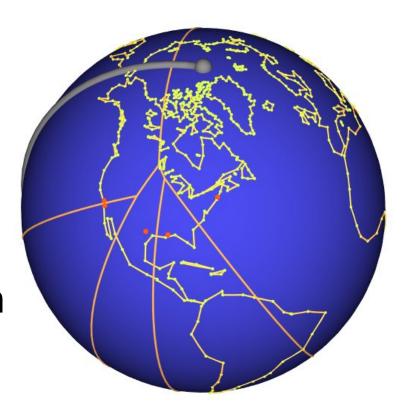
Simultaneous Polyline Simplification

- Implementation of [Dyken et al]
- Based on CGAL::Constrained_delaunay_2
- Guarantees that after simplification
 - islands stay islands
 - isolines do not intersect



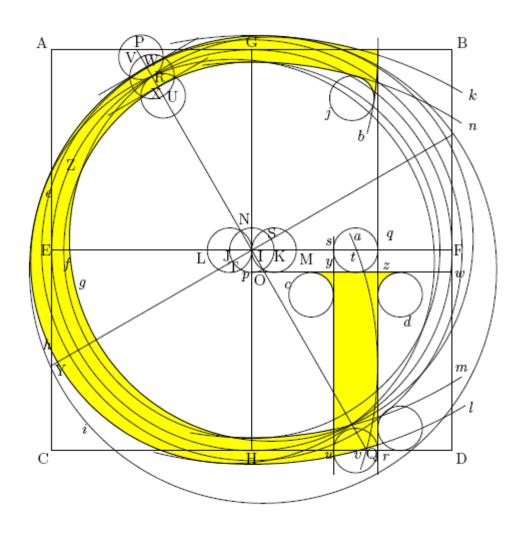
Vector Graphics on the Sphere

- Arrangement_2<Geometry, Embedding>
- Boolean operations
- Map overlay
- Voronoi diagram
- Point location
- Convex decomposition



Summary: CGAL for Vector Graphics

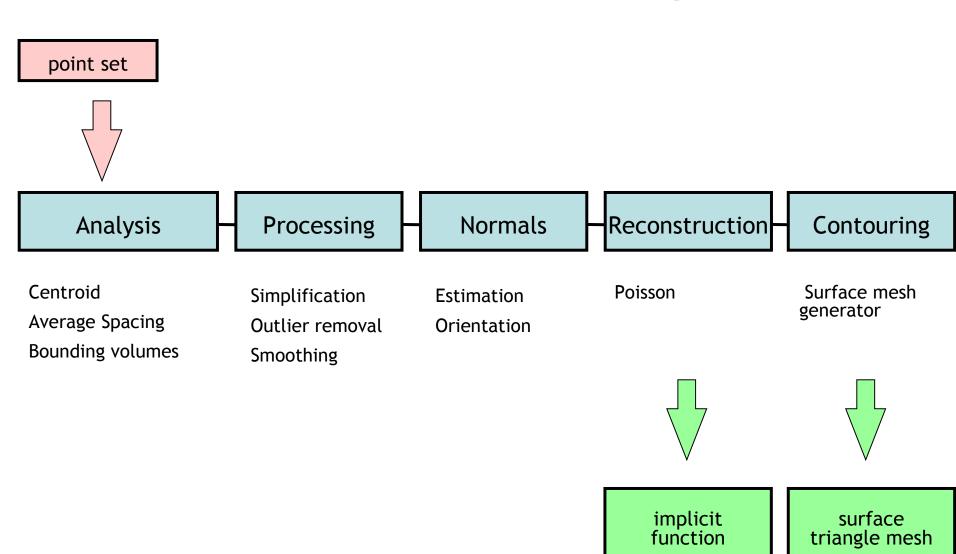
- Rich collection of 2D geometric algorithms
- Modular and generic design
- Linear and curved primitives
- Useful in many application domains



CGAL for Point Sets

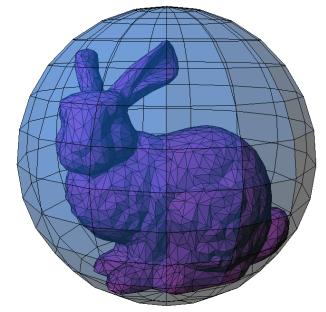
Pierre Alliez INRIA

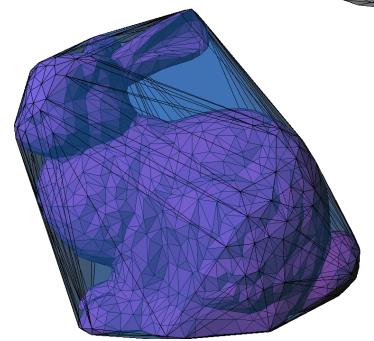
Surface Reconstruction Pipeline



Bounding Volumes

- Convex hull
- Bounding sphere





Principal Component Analysis

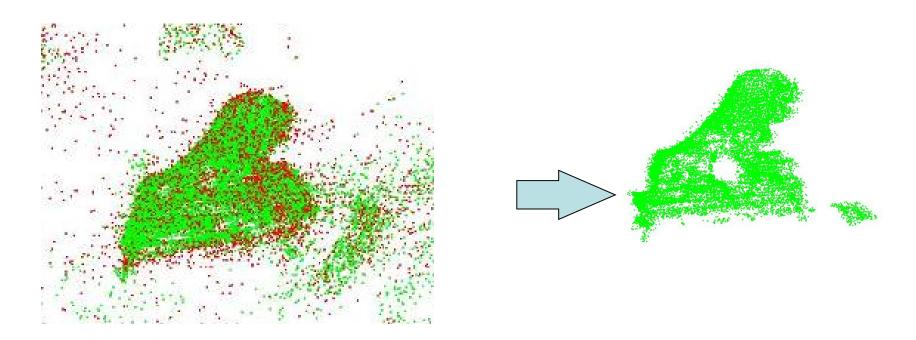
Linear least squares fitting on sets of 3D points



CGAL manual

Outlier Removal

 Sort w.r.t. sum of squared distances to k-nearest neighbors (CGAL::K_nearest_neighbor_search) and cut at specified percentile.



Estimation of Curvatures

- Estimates general differential properties (Monge form) on point sets.
- Through polynomial (d-jet) fitting



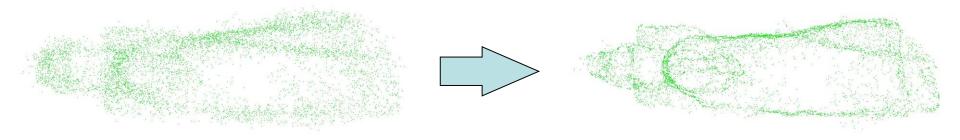
min curvature directions



max curvature directions

Point Cloud Smoothing

- For each point
 - Find k-nearest neighbors
 - fit jet (smooth parametric surface)
 - project onto jet



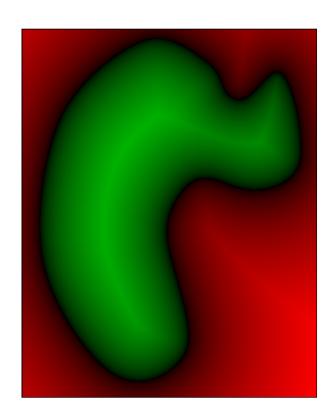
(noisy point set)

(smoothed point set)

Surface Reconstruction

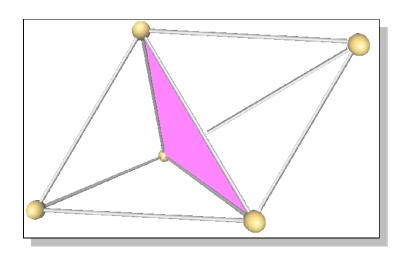
Poisson Surface Reconstruction [Kazhdan-Bolitho-Hoppe, SGP 2006]

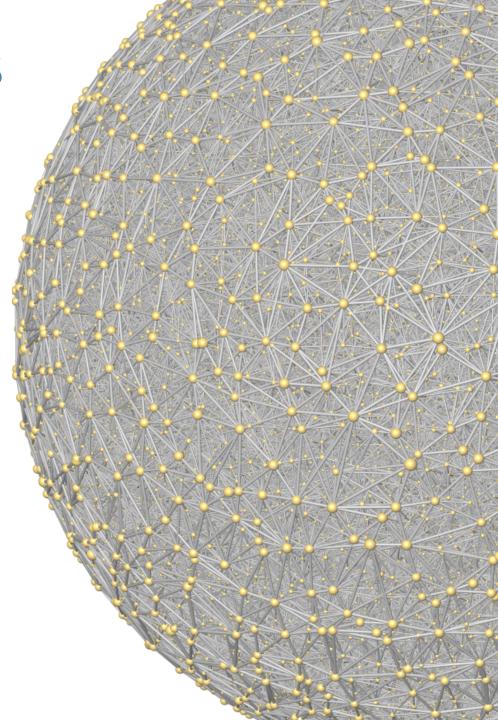
- Solves for an implicit function (~indicator function)
- Isosurface extracted by CGAL::Surface_mesher



3D Triangulations

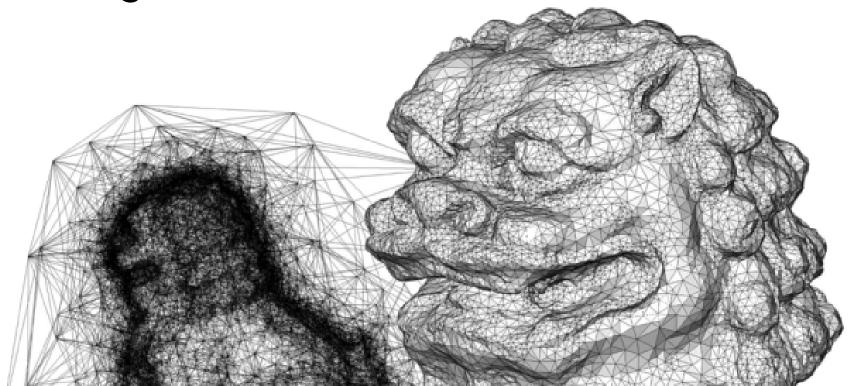
- Delaunay
- Fully dynamic
- 1M 3D points in 16 sec





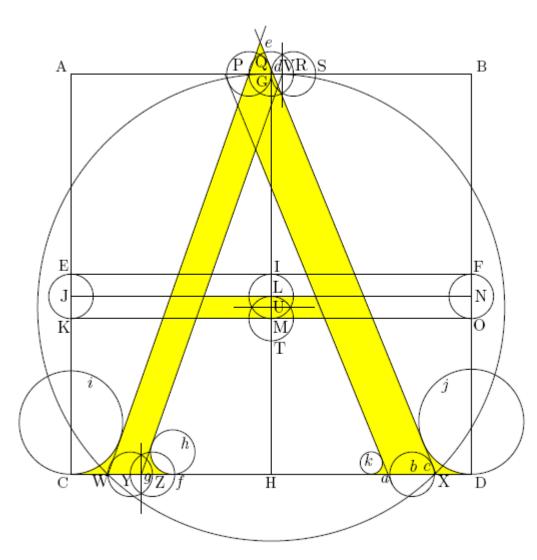
Surface Reconstruction Demo

 Solves for the Poisson equation onto the vertices of a (refined) 3D Delaunay triangulation.



Summary: CGAL for Point Sets

- Algorithms are modular components
 - in this course: positioned along the surface reconstruction pipeline.
 - can be used individually
- Poisson reconstruction is the first algorithm of the surface reconstruction package.



CGAL for Modeling and Processing of Polyhedral Surfaces

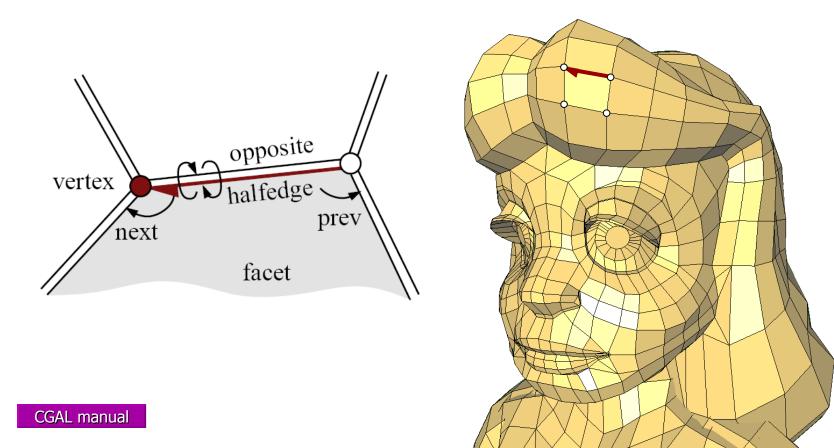
Andreas Fabri GeometryFactory

Outline

- Polyhedral Surface
 - Halfedge data structure
 - Euler Operators
 - Customization
- Algorithms for Geometric Modelling and Geometry Processing

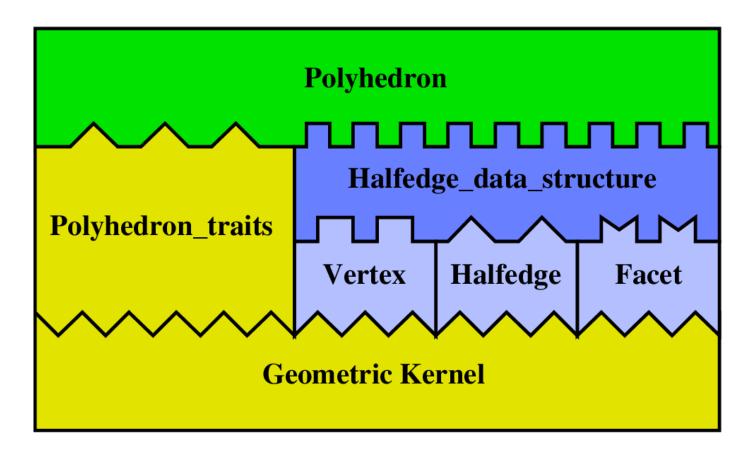
Halfedge Data Structure

Represented by vertices, edges, facets and an **incidence relation** on them, restricted to orientable 2-manifolds with boundary.



Polyhedron

Building blocks assembled with C++ templates

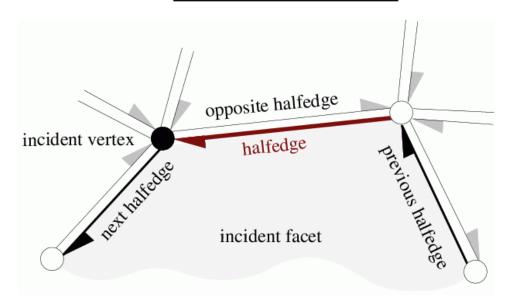


Default Polyhedron

Vertex Halfedge_handle halfedge() Point& point()

Halfedge	
Halfedge_handle	opposite()
Halfedge_handle	next()
Halfedge_handle	prev()
Vertex_handle	vertex()
Facet_handle	facet()

Facet	
Halfedge_handle	halfedge()
Plane&	plane()
Normal&	normal()
Color&	color()

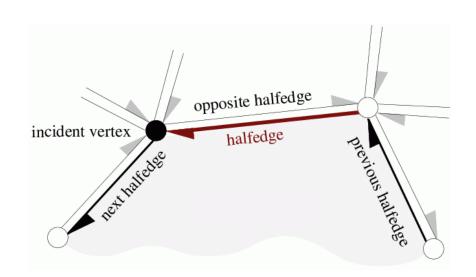


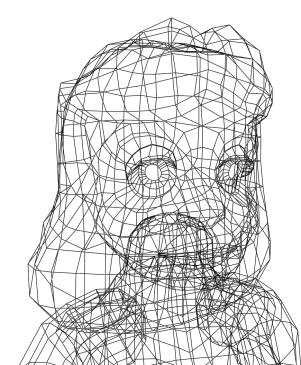
Flexible Data Structure

Vertex Halfedge_handle halfedge() Point& point()

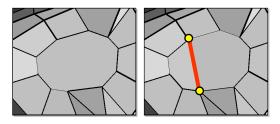
Halfedge	
Halfedge_handle	opposite()
Halfedge_handle	next()
Halfedge_handle	prev()
Vertex_handle	vertex()
Facet_handle	facet()

Facet	
Halfedge_handle	halfedge()
Plane&	plane()
Normal&	normal()
Color&	color()

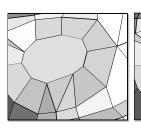




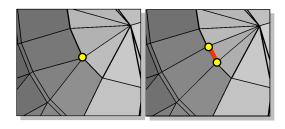
Euler Operators



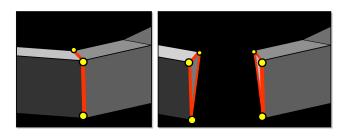
split_facet
join_facet



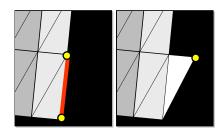
create_center_vertex
erase_center_vertex



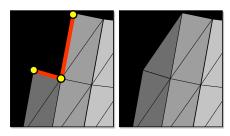
split_vertex join_vertex (aka edge collapse)



split_loop
join_loop



add_vertex_and_facet _to_border erase_facet



add_facet_to_border erase_facet

CGAL manual

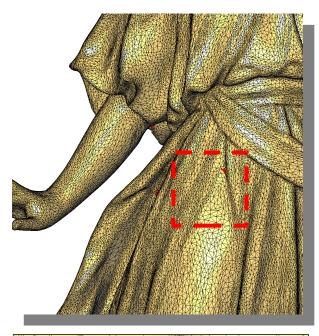
Algorithms

Algorithms

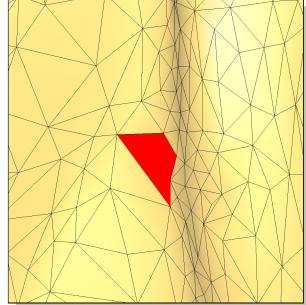
- Intersection detection
- AABB Tree
- Bounding volumes
- Boolean operations
- Kernel
- Parameterization
- Subdivision
- Principal component analysis
- Extraction of ridges
- Simplification

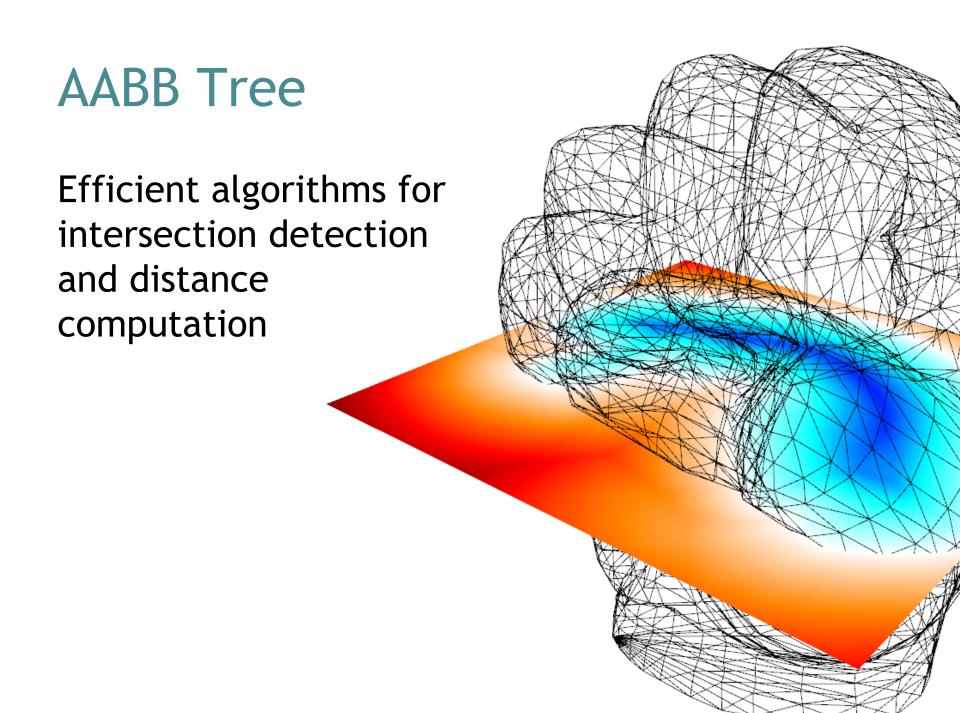
Intersection Detection

- Efficient algorithm for finding all intersecting pairs for large numbers of axis-aligned bounding boxes.
- Generic programming:
 Boxes can contain
 objects of any type







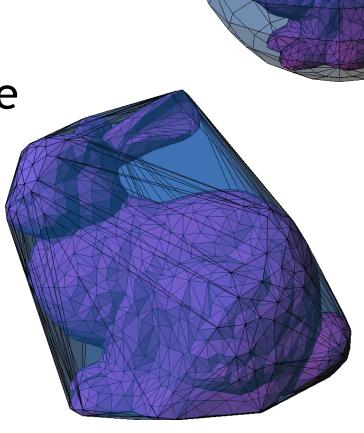


Bounding Volumes

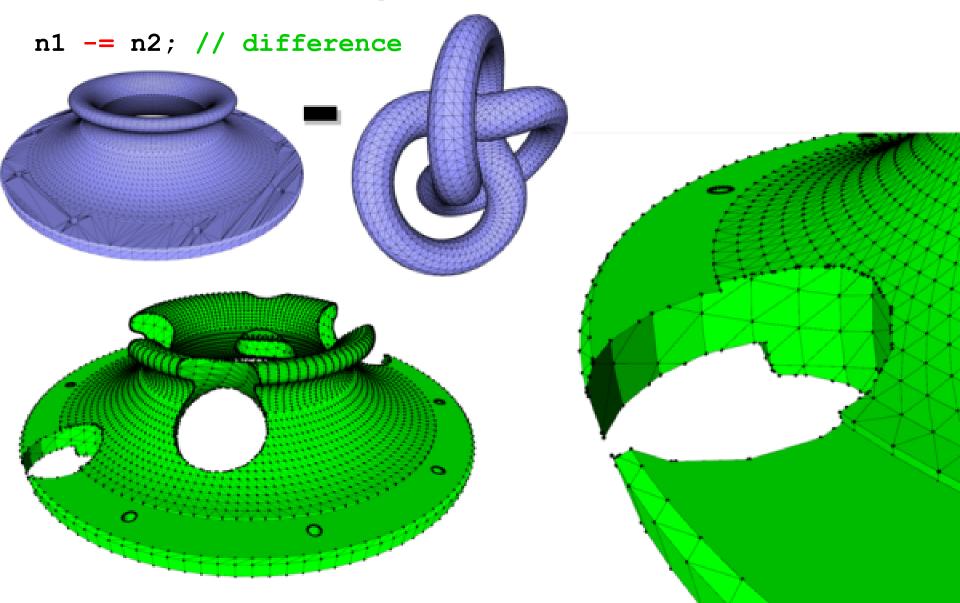
Convex hull

Bounding sphere

Bounding sphere
 of spheres

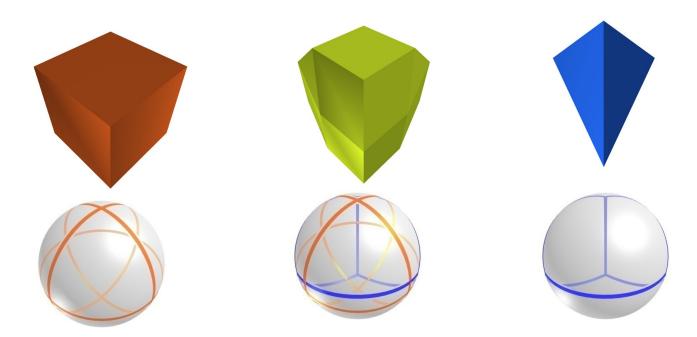


3D Boolean Operations



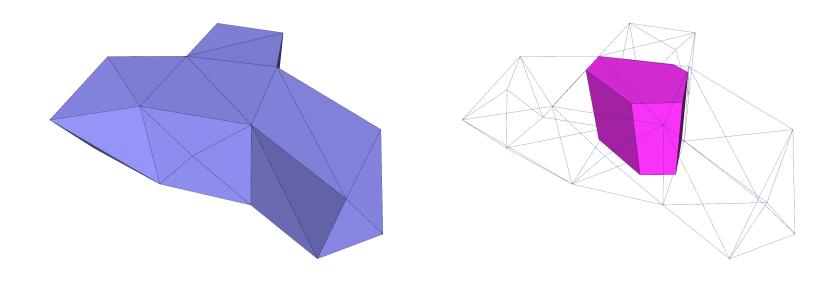
Minkowski-Sums of Polytopes

- The Gaussian map of a polytope is the decomposition of S2 into maximal connected regions so that the extremal point is the same for all directions within one region
- The overlay of the Gaussian maps of two polytopes P and Q is the Gaussian map of the Minkowski sum of P and Q



Kernel of a Polyhedron

- Intersection of all its interior half-spaces
- Uses linear programming: CGAL::QP_solver



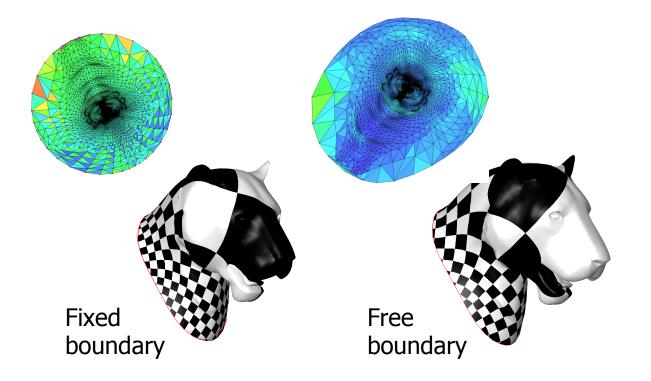
input polyhedron

kernel

Parameterization

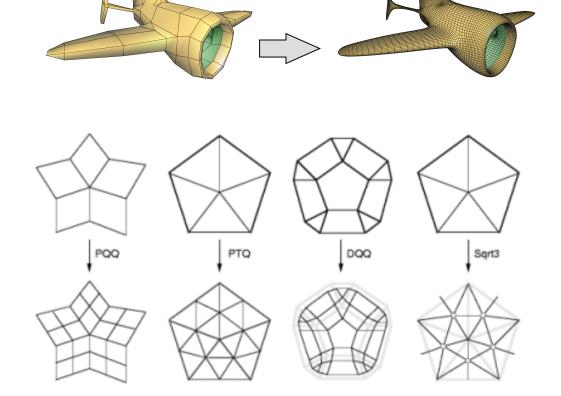
- Planar
- Conformal [Eck et al., Levy et al., Desbrun et al.]
- Mean value coordinates [Floater]

• ...



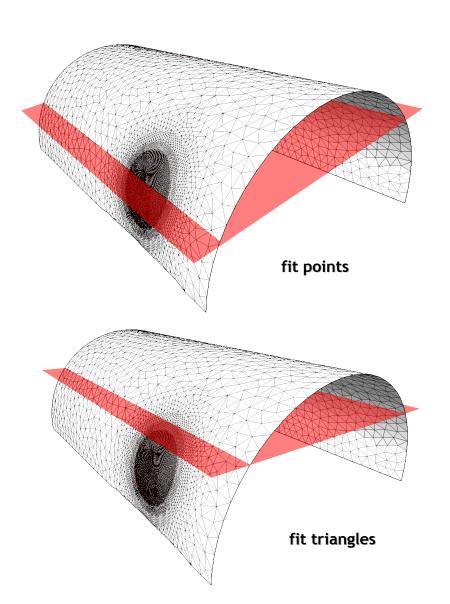
Subdivision

- Designed to work on CGAL polyhedron
- Catmull-Clark
- Loop
- Doo-Sabin
- Sqrt3
- ...



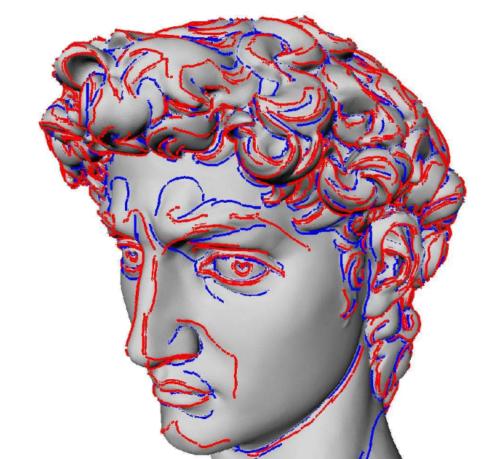
Principal Component Analysis

 Linear least squares fitting on sets of 3D points or triangles



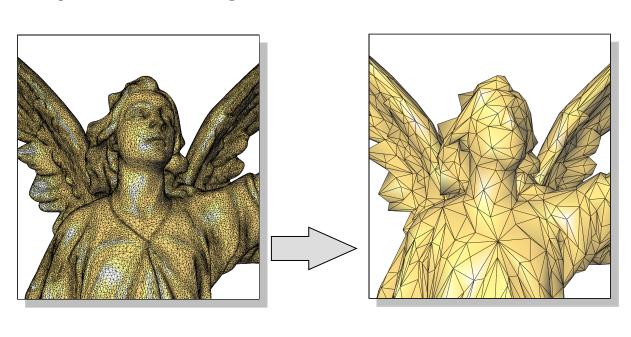
Extraction of Ridges

 Ridge: curve along which one of the principal curvatures has an extremum along its curvature line.



Simplification

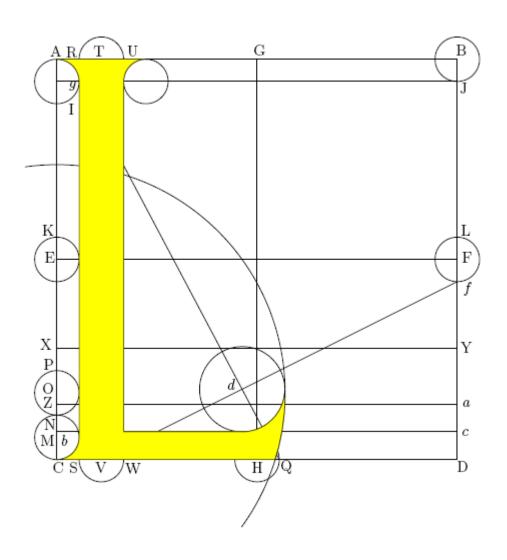
 Implementation of [Lindstrom-Turk] volumepreserving method.





Summary: CGAL for Modeling and Polyhedral Surfaces

- The halfedge data structure and the polyhedron are highly flexible
- CGAL provides many algorithms for geometric modeling and geometry processing
- Polyhedral surface as output of surface mesh generation algorithms



CGAL for Mesh Generation

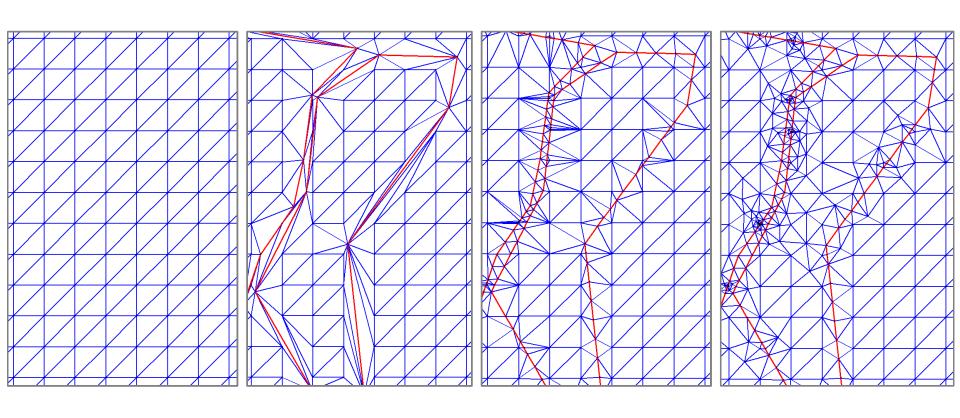
Pierre Alliez INRIA

Outline

- 2D mesh generation
- Surface mesh generation
- 3D mesh generation
- Work in progress

2D Mesh Generation

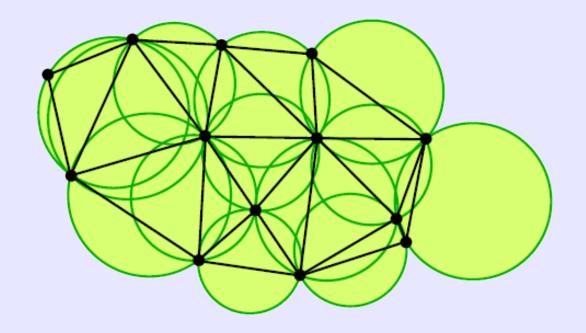
2D Mesh Generation



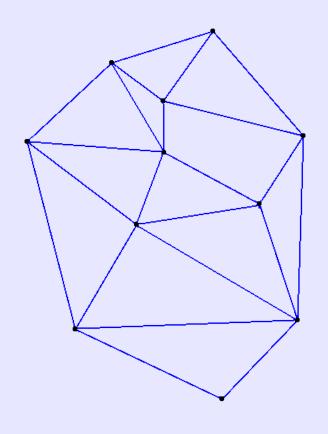
From Triangulations to Quality Meshes

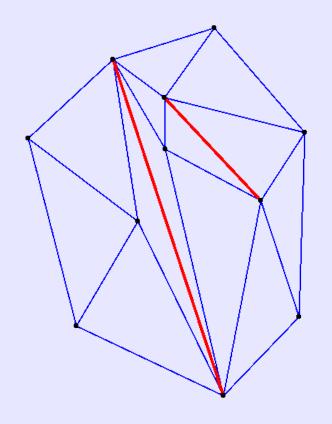
Delaunay Triangulation

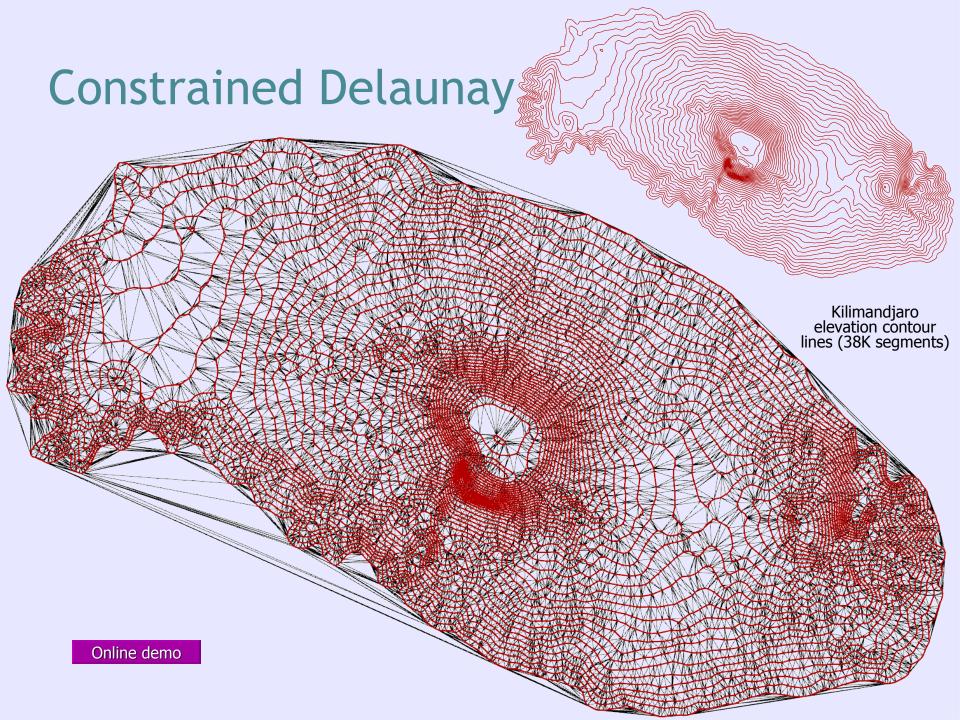
 A triangulation is a Delaunay triangulation, if the circumscribing circle of any facet of the triangulation contains no vertex in its interior



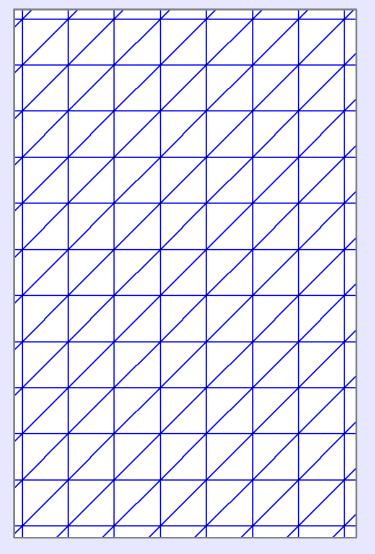
Constrained Delaunay Triangulation

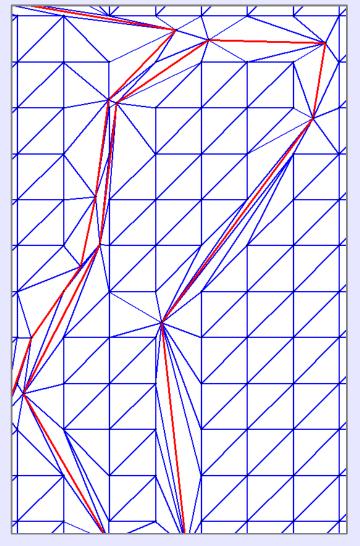




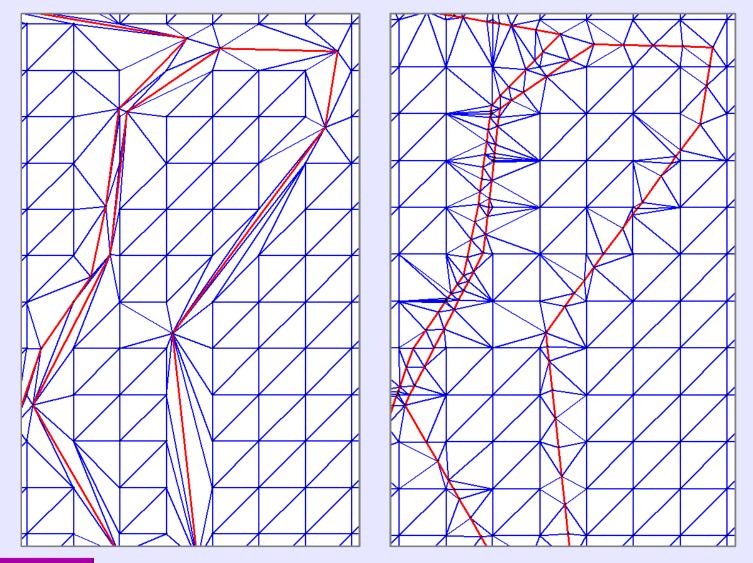


Adding Constraints

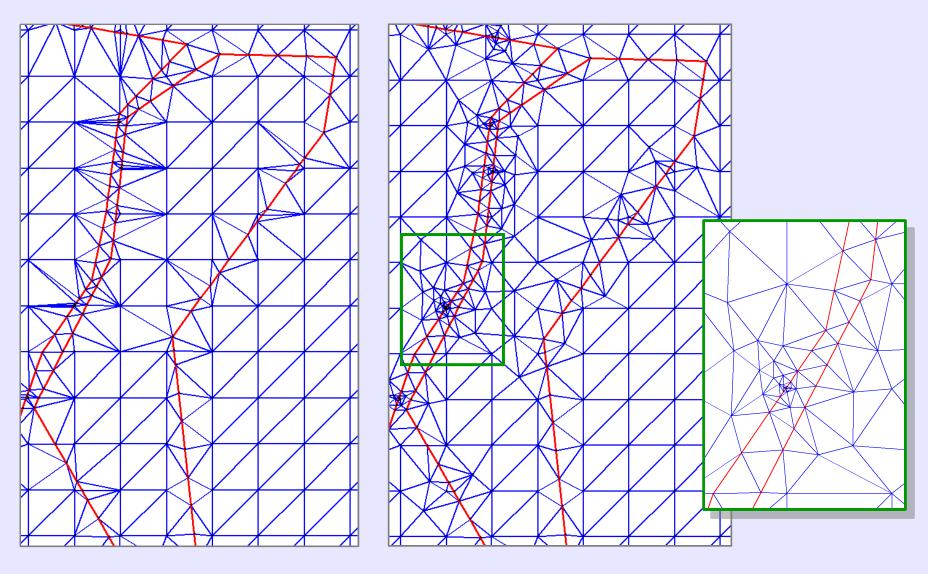




Conforming a Triangulation

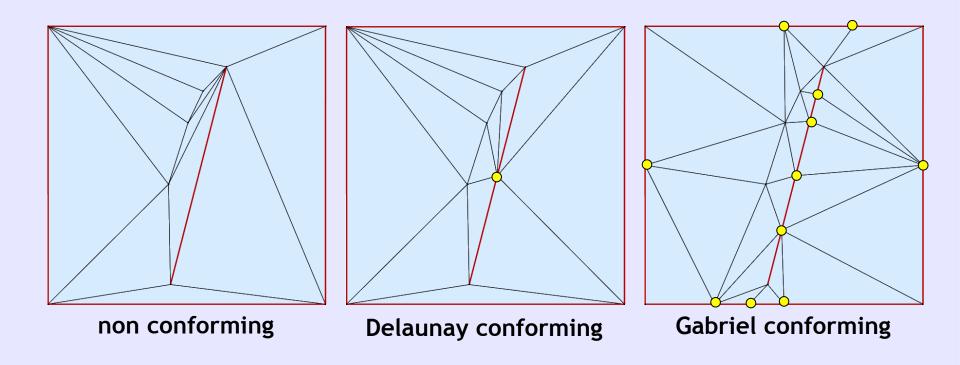


Delaunay Refinement



Conforming a Triangulation

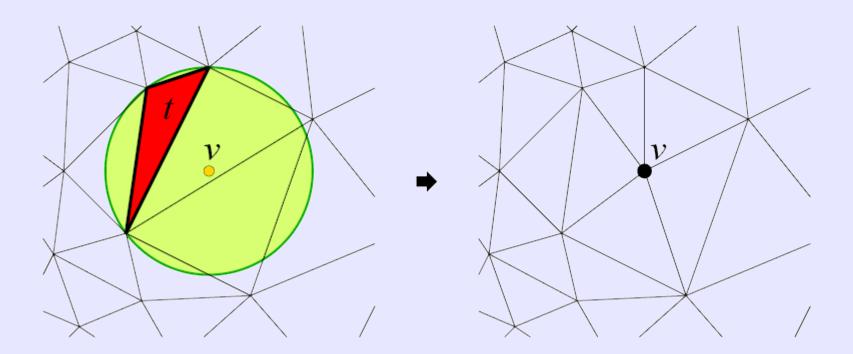
Any constrained Delaunay triangulation can be refined into a conforming Delaunay or Gabriel triangulation by adding Steiner vertices.



Delaunay Refinement Rules

Rule #1: break bad elements by inserting circumcenters (Voronoi vertices)

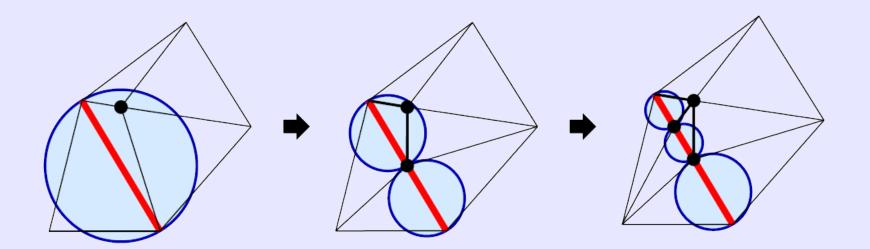
- "bad" in terms of size or shape (too big or skinny)



Delaunay Refinement Rules

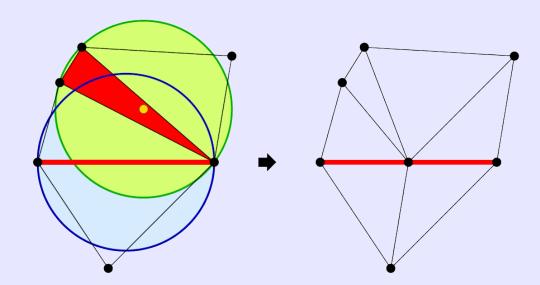
Rule #2: Midpoint vertex insertion

A constrained segment is said to be encroached, if there is a vertex inside its diametral circle



Delaunay Refinement Rules

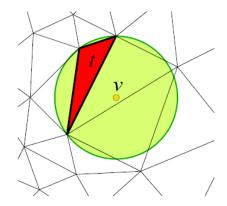
Encroached subsegments have priority over skinny triangles

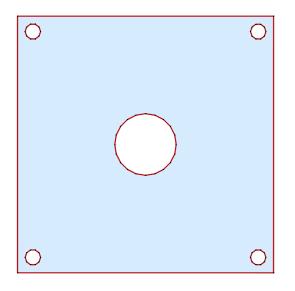


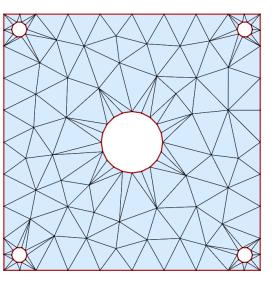
Parameters for Delaunay Refinement

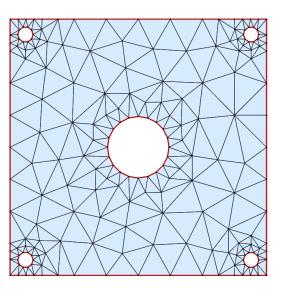
Shape

- Lower bound on triangle angles









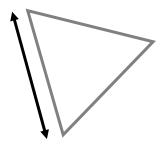
Input PLSG

5 deg

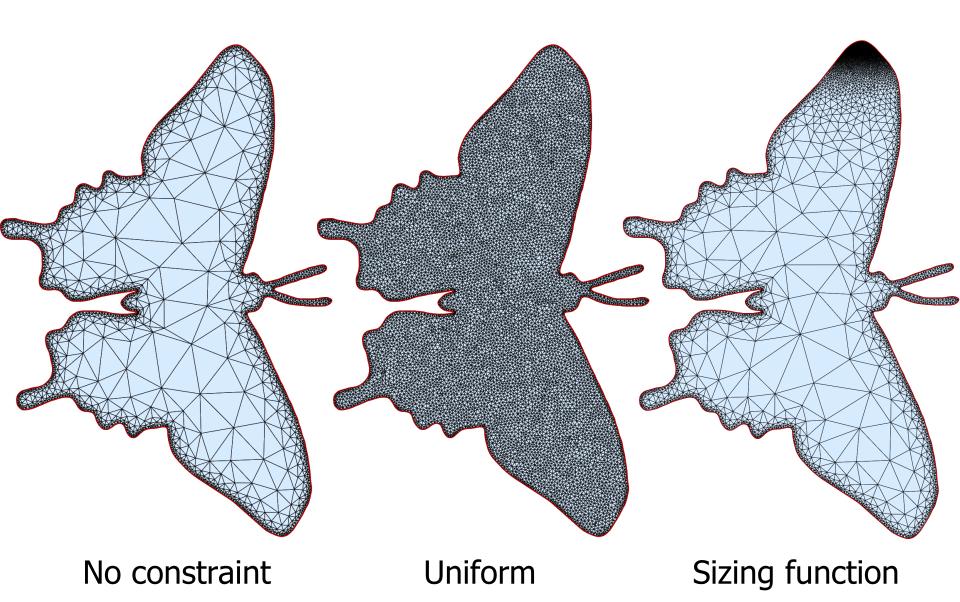
20.7 deg

Parameters for Delaunay Refinement

- Shape
 - Lower bound on triangle angles
- Size
 - No constraint
 - Uniform sizing
 - Sizing function

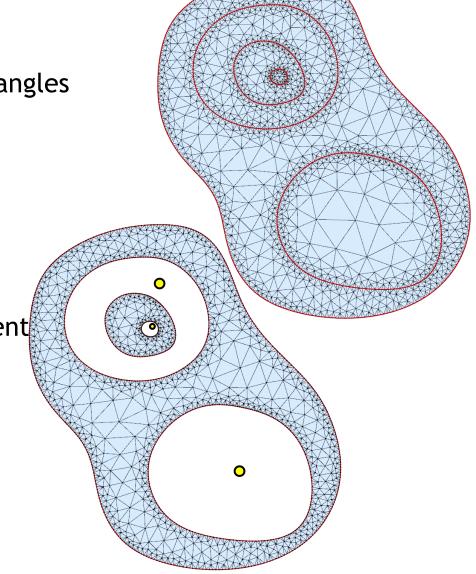


Sizing Parameter

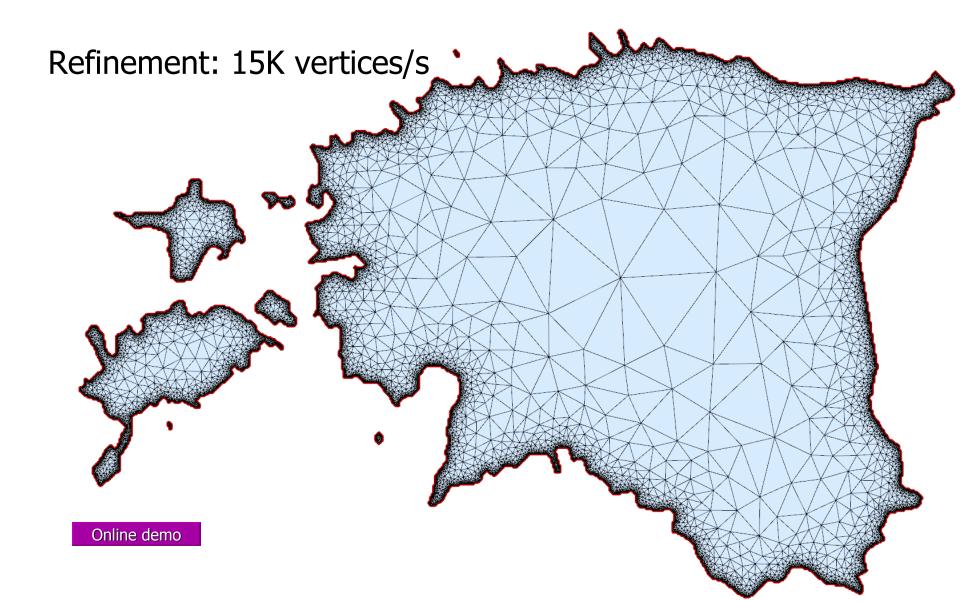


Parameters for Delaunay Refinement

- Shape
 - Lower bound on triangle angles
- Size
 - No constraint
 - Uniform sizing
 - Sizing function
- Seeds
 - Exclude/include component

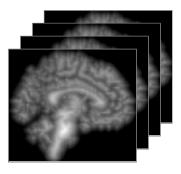


Performances

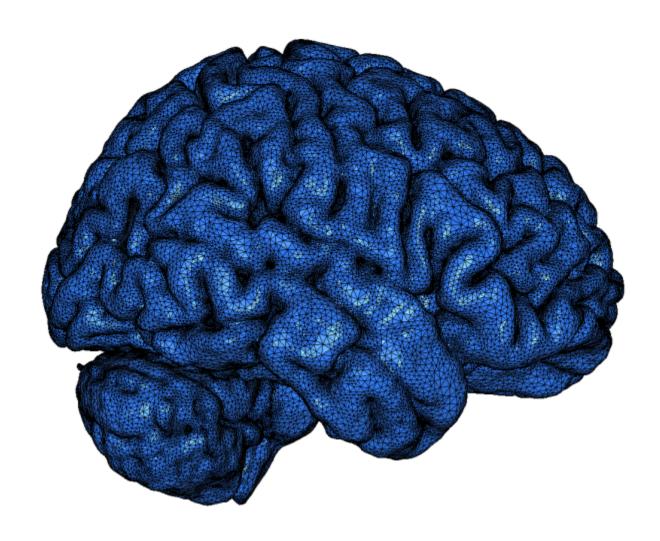


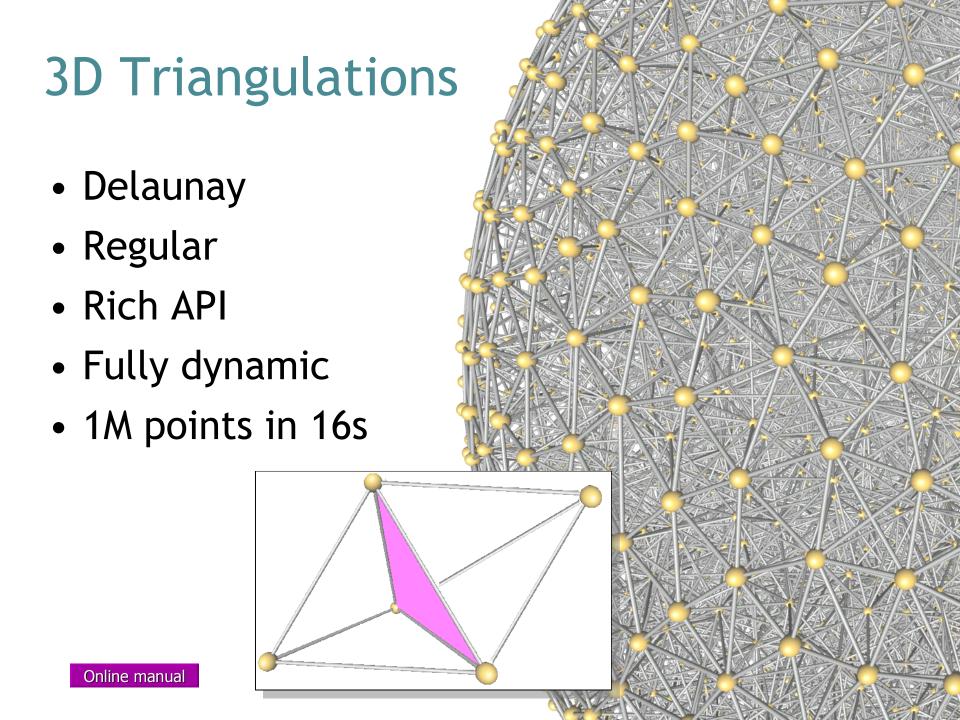
Surface Mesh Generation

Surface Mesh Generation



input

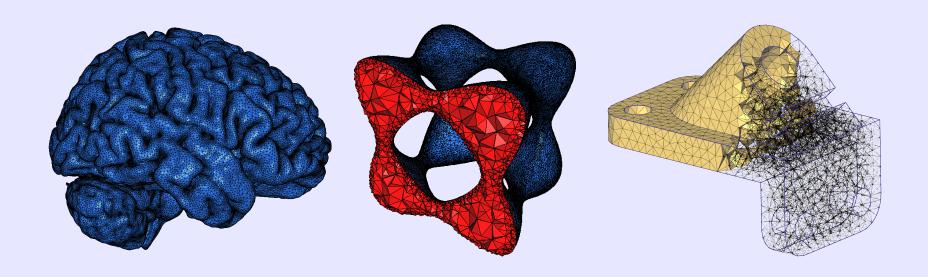




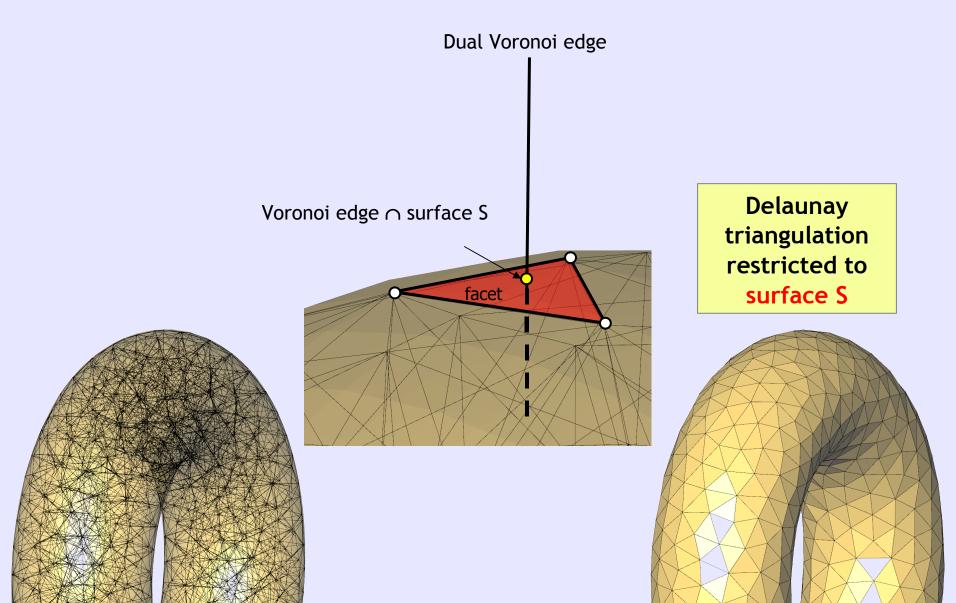
Mesh Generation

Key concepts:

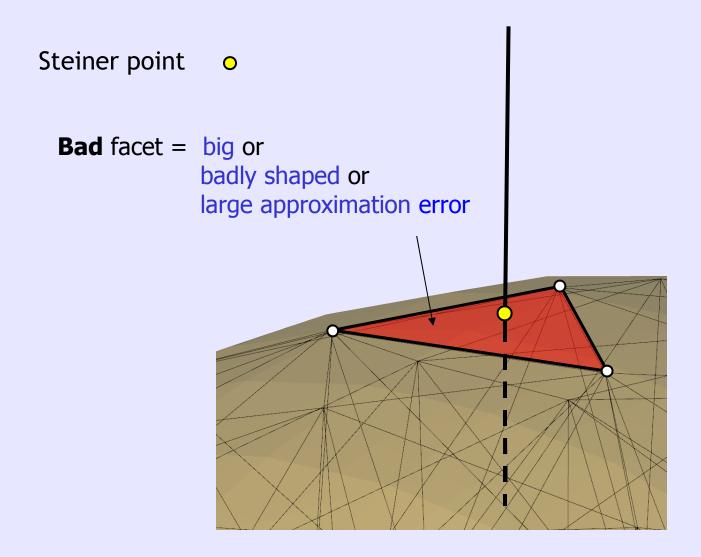
- Delaunay filtering
- Delaunay refinement



Delaunay Filtering



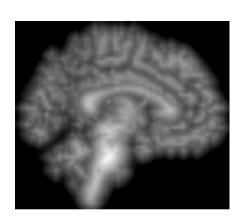
Delaunay Refinement



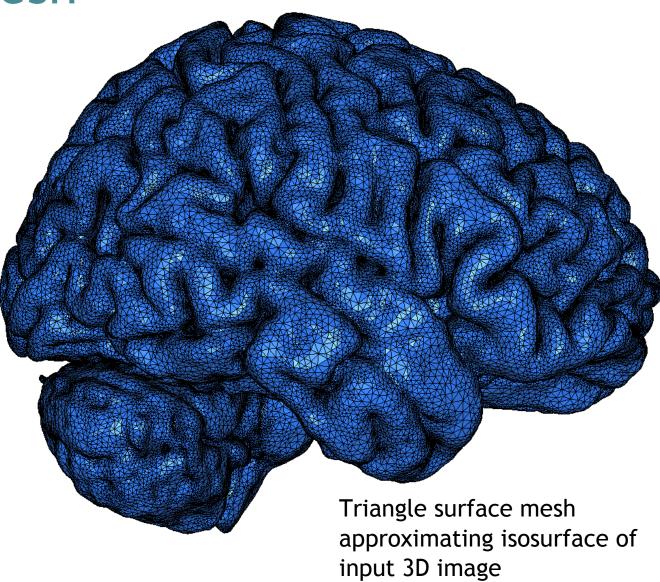
Delaunay Refinement

```
repeat \{ pick bad facet f insert furthest (dual(f) \cap S) in Delaunay triangulation update Delaunay triangulation restricted to S \} until all facets are good
```

Output Mesh



input



Output Mesh

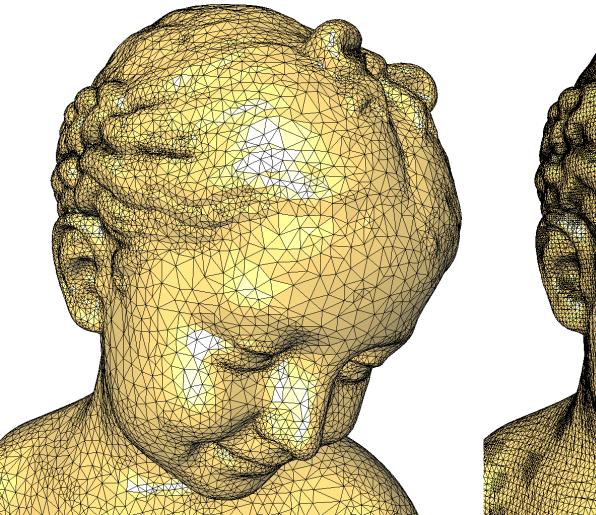
- Well shaped triangles
 - Lower bound on triangle angles
- Homeomorphic to input surface
- Manifold
 - not only combinatorially, i.e., no self-intersection
- Faithful Approximation of input surface
 - Hausdorff distance
 - Normals

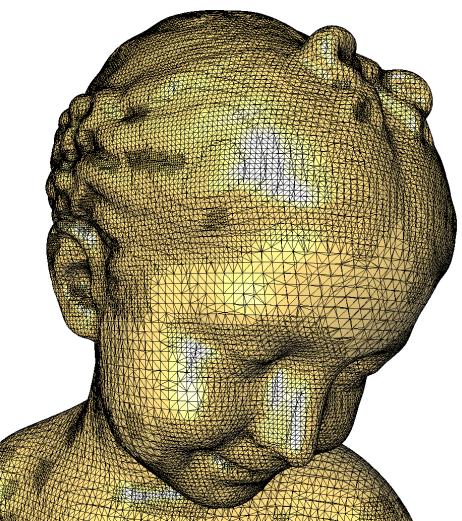


vs Marching Cubes

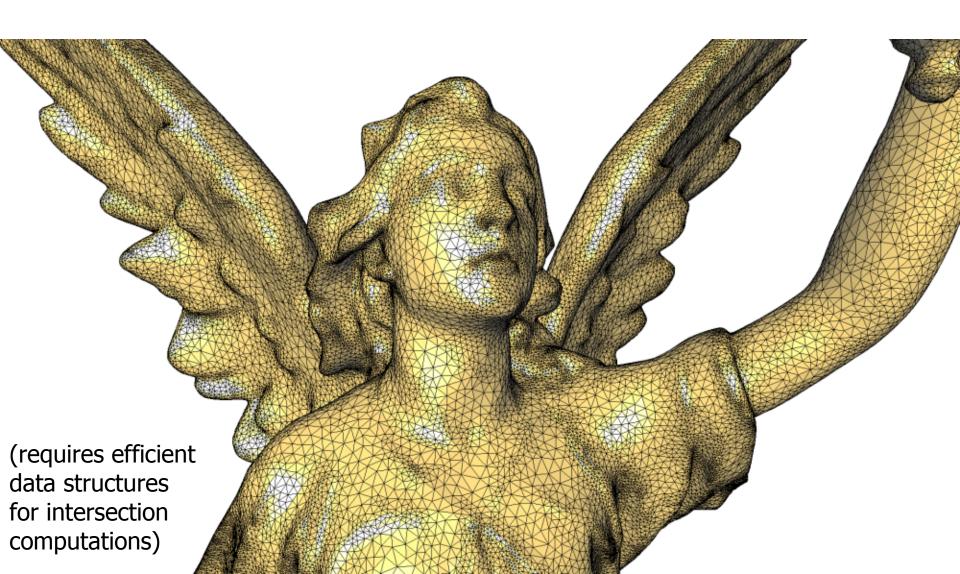
Delaunay refinement

Marching cubes in octree





Surface Remeshing (input is a polyhedral surface)

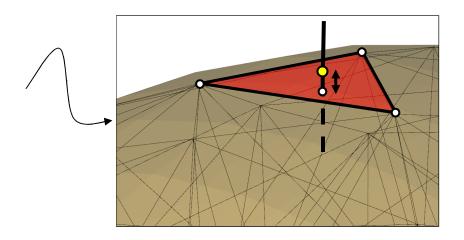


Parameters

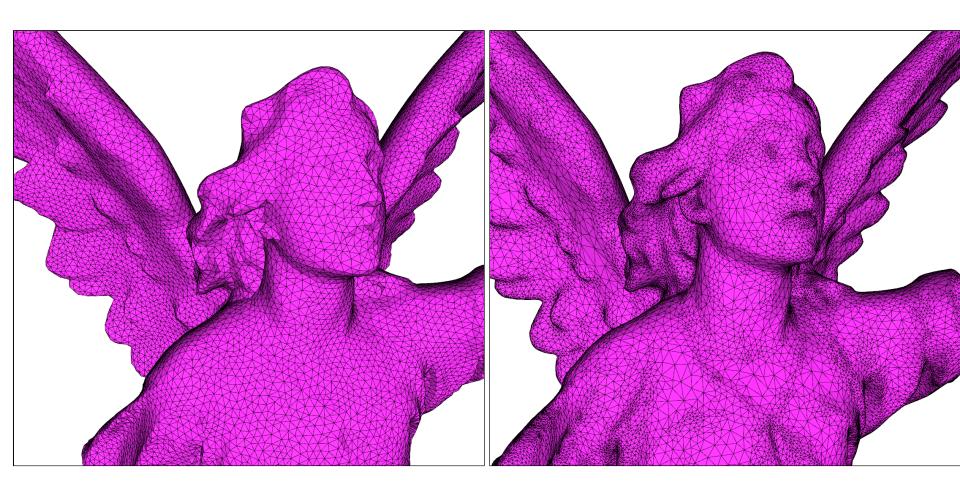
- Shape of triangles
 - lower bound on triangle angles
- Size
 - No constraint
 - Uniform sizing
 - Sizing function

Parameters

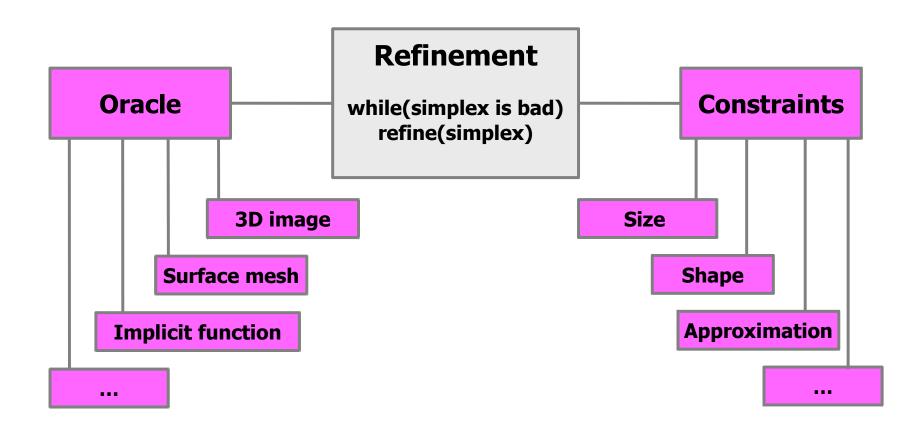
- Shape of triangles
 - lower bound on triangle angles
- Size
 - No constraint
 - Uniform sizing
 - Sizing function
- Approximation error



Uniform vs Adapted



Mesh Generation Framework

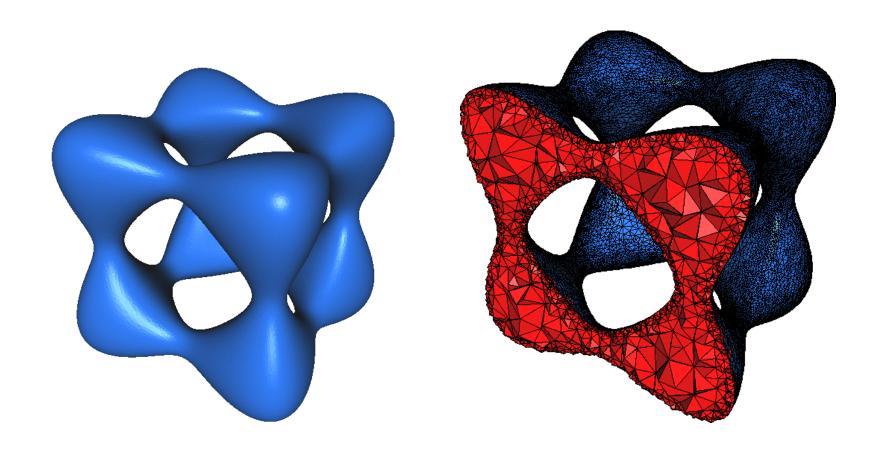


A Versatile Framework

- 3D grey level images
- 3D multi-domain images
- Implicit function: f(x, y, z) = constant
- Surface mesh (remeshing)
- Point set (surface reconstruction)
- Anything which provides intersections

3D Mesh Generation

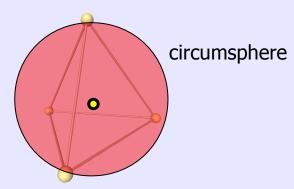
3D (Volume) Mesh Generation



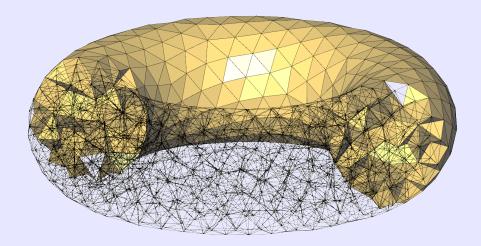
More Delaunay Filtering

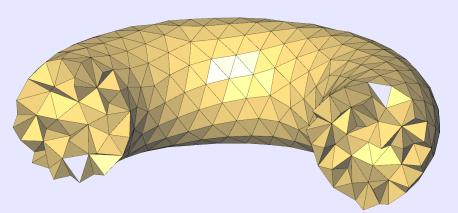
Delaunay triangulation restricted to domain Ω

tetrahedron



Dual Voronoi vertex inside domain Ω ("oracle")

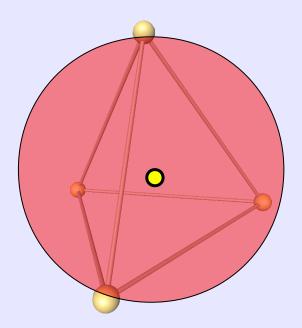




Delaunay Refinement

Steiner point

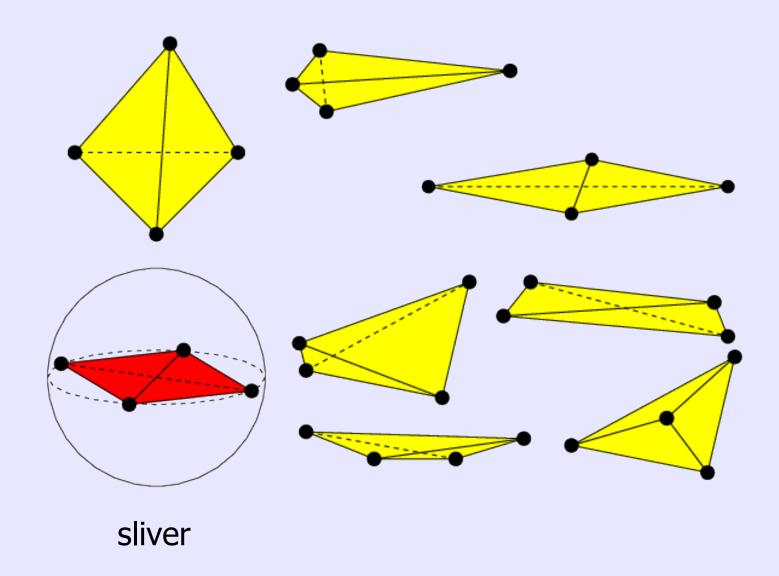
Bad tetrahedron = big or badly shaped



Volume Mesh Generation Algorithm

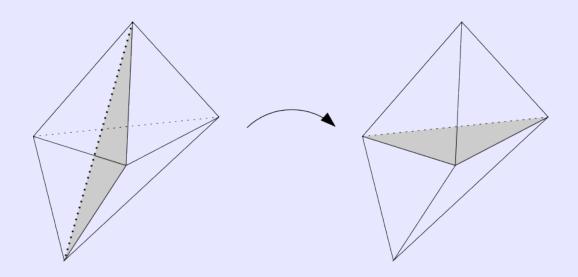
```
repeat
pick bad simplex
if(Steiner point encroaches a facet)
  refine facet
else
  refine simplex
update Delaunay triangulation restricted to domain
until all simplices are good
Exude slivers
```

Tetrahedron Zoo



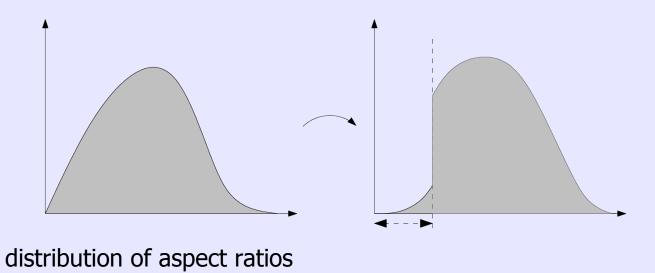
Sliver Exudation [Edelsbrunner-Guoy]

- Delaunay triangulation turned into a regular triangulation with null weights.
- Small increase of weights triggers edge-facets flips to remove slivers.

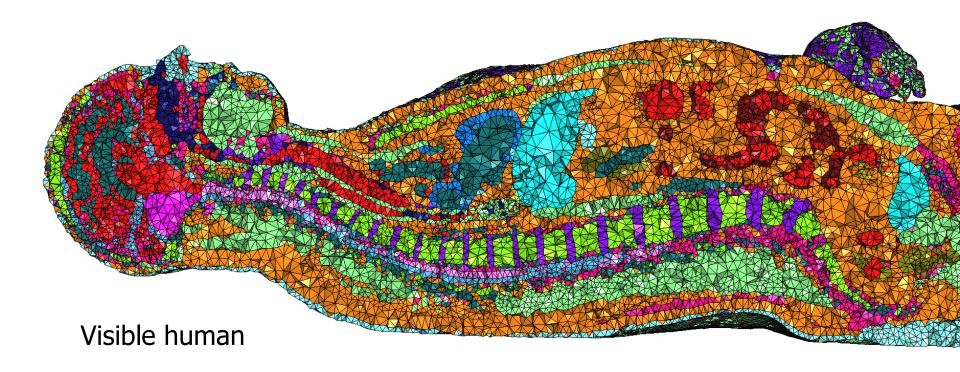


Sliver Exudation Process

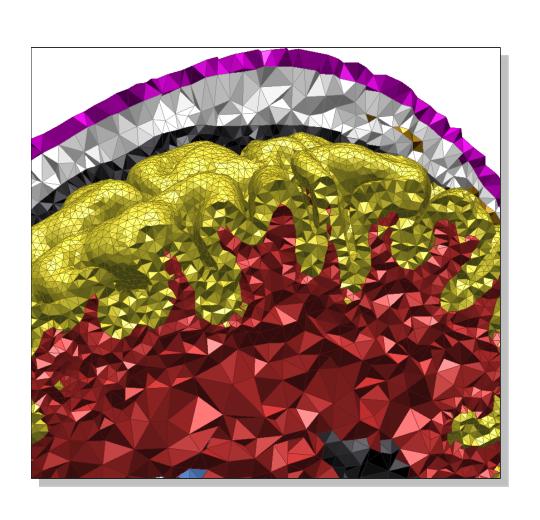
- Try improving all tetrahedra with an aspect ratio lower than a given bound
- Never flips a boundary facet

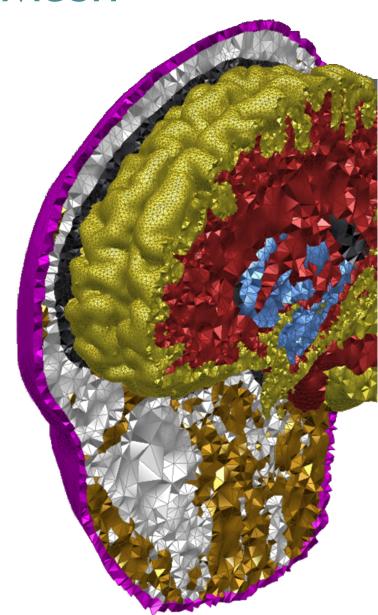


Multi-Domain Volume Mesh



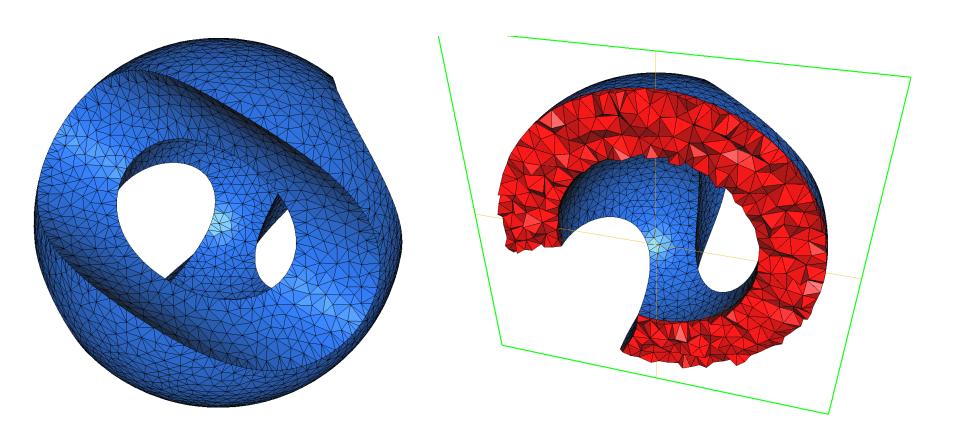
Multi-Domain Volume Mesh



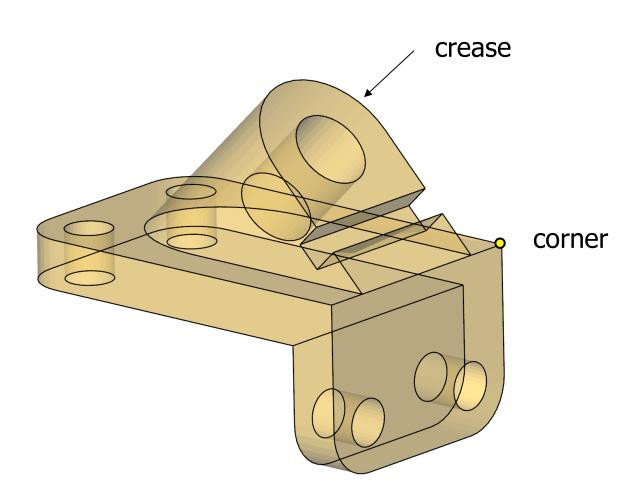


Work in Progress

Piecewise Smooth Surfaces



Input: Piecewise smooth complex



Even More Delaunay Filtering

primitive

dual of

Voronoi vertex

tetrahedron

Voronoi edge

facet

Voronoi face

edge

test

against

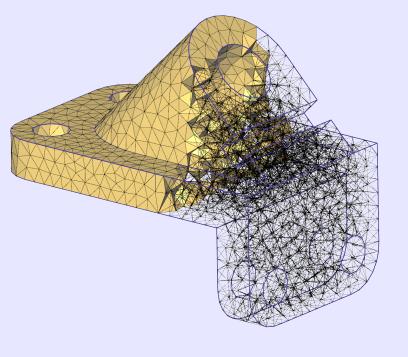
inside

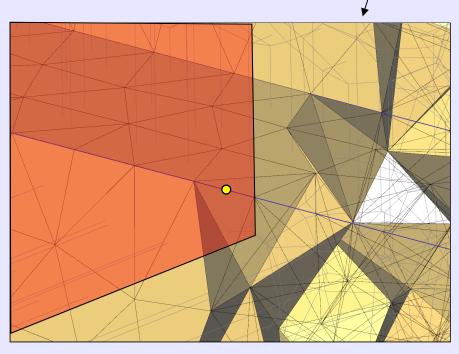
domain

intersect

domain boundary

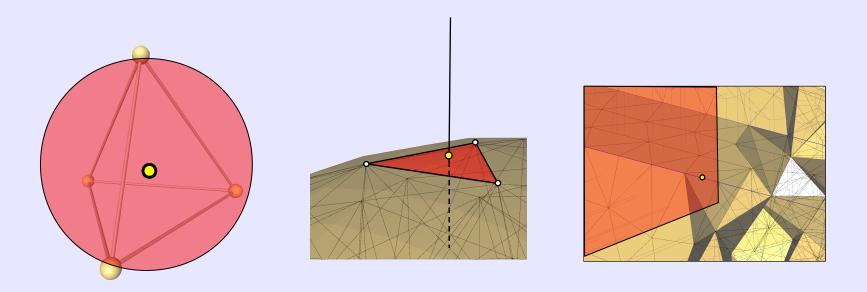
intersect crease





Delaunay Refinement

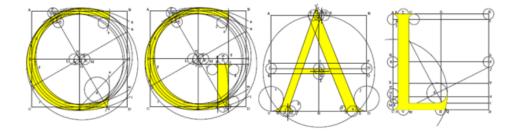
Steiner points



Summary: CGAL for Mesh Generation

- 2D mesh generation
 - From triangulation to quality mesh
 - Preserves constraints exactly

- 3D Mesh generation
 - Interpolates boundary
 - Versatile through oracle-based design



Questions and Answers

Andreas Fabri GeometryFactory

Pierre Alliez INRIA

Question and Answers

- General Introduction
- CGAL for 2D Vector Graphics
- CGAL for Point Sets
- CGAL for Modeling and Processing of Polyhedral Surfaces
- CGAL for Mesh Generation